

Lecture 16 (Planar graph)

A graph is said to be embedded on a surface S when it is drawn on S so that no two edges intersect, except at endpoints. A graph is said to be planar if it can be embedded on the plane.

Examples:

1. A tree is embeddable on a plane.
2. Any cycle C_n , $n \geq 3$ is planar.
3. The K_4 is planar.
4. The $K_{2,3}$ is planar.
5. Draw a planar embedding of $K_5 - e$, where e is any edge.
6. Draw a planar embedding of $K_{3,3} - e$, where e is any edge.

Definition 1. Consider a planar embedding of a graph G . The regions on the plane defined by this embedding are called faces/regions of G . The unbounded face/region is called the exterior face.

Theorem 0.1. (Euler Formula for Planar Graph) Let G be a connected planar graph with V number of vertices, E number of edges and F number of faces. Then $V - E + F = 2$.

Proof: We use induction on E . Let $E = 1$. Then $V = 2$ and $F = 1$ and hence $V - E + F = 2$.

Suppose the result holds for all the planar-connected graphs with $E \leq n - 1$.

Let G be a connected planar graph with $E = n$. Then we have two cases:

Case 1: If G has no cycles. Then G is a tree and hence the result is true.

Case 2: Suppose G contains at least one cycle. Let e be an edge in a cycle and $G' = G - e$. Then $E' = E - 1$, $V' = V$ and $F' = F - 1$. Then $E' - V' + F' = 2$ implies $V - E + F = 2$. Thus the result holds.

Corollary 1. If G is a connected planar simple graph with E edges and V vertices, where $V \geq 3$, then $E \leq 3V - 6$.

Proof: Each face has at least 3 edges and each edge is participated in two faces. So, the number of edges $E \geq 3F/2$. Now the proof follows by $V - E + F = 2$.

Exercise: K_5 is non-planar, because here $V = 5$ and $E = 10$, using the above result leads to $E = 10 \leq 3 \times 5 - 6 = 9$, a contradiction.

Corollary 2. If a connected planar simple graph has E edges and V vertices with $V \geq 3$ and no cycle of length three, then $E \leq 2V - 4$.

Proof: Note that $E \geq 2F$. Then the proof follows by $V - E + F = 2$.

Exercise: $K_{3,3}$ is non-planar, because $K_{3,3}$ has no cycle of length 3 with $V = 6, E = 9$, which using the above result leads to $E = 9 \leq 2 \times 6 - 4 = 8$, a contradiction.

Proof: Note that $E \geq 2F$. Then the proof follows by $V - E + F = 2$.

Remark. The following result generalizes the above corollaries for a connected, planar graph having cycles of any length.

Theorem 0.2. Let G be a connected planar graph with E edges and V vertices, $V \geq 3$, such that G has cycles of length $\geq g$, then $E \leq \frac{g}{g-2}(V - 2)$.

Proof: Suppose the embedding has F faces. Then $E \geq \frac{gF}{2} \Leftrightarrow F \leq \frac{2E}{g}$. Now following Euler formula, we get $E \leq \frac{g}{g-2}(V - 2)$.

Definition 2. Let G be a graph. Then, a subdivision of an edge uv in G is obtained by replacing the edge by two edges uw and wv , by introducing a new vertex w of degree 2. Two graphs G_1 and G_2 are said to be homeomorphic if there is a graph G_3 such that both graphs G_1 and G_2 can be obtained from G_3 a sequence of subdivisions.

Example. For $n \neq m$, the path graphs P_n is homeomorphic to P_m , as P_n and P_m can be obtained from P_1 by applying $n - 2$ and $m - 2$ numbers of subdivisions respectively. Similarly C_n is homeomorphic to C_m for distinct m and n . However, C_n is not homeomorphic to wheel graph W_n , where $V(W_n) = V(C_n) \cup \{a\} = \{v_1, v_2, \dots, v_n\} \cup \{a\}$, and $E(W_n) = E(C_n) \cup \{av_1, av_2, \dots, av_n\}$.

Theorem 0.3. (Kuratowski's Theorem) A graph is non planar if and only if it contains a subgraph homeomorphic to $K_{3,3}$ or K_5 .

Exercise: Petersen graph is non-planar. It can be shown that the said graph is homeomorphic to $K_{3,3}/K_5$.